The design of a computer language for linguistic information

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The Design of a Computer Language for Linguistic Information
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Abstract

A considerable body of accumulated knowledge about the design of languages for communicating information to computers has been derived from the subfields of programming language design and semantics. It has been the goal of the PATR group at SRI to utilize a relevant portion of this knowledge in implementing tools to facilitate communication of linguistic information to computers. The PATR-II formalism is our current computer language for encoding linguistic information. This paper, a brief overview of that formalism, attempts to explicate our design decisions in terms of a set of properties that effective computer languages should incorporate.

1. Introduction

The goal of natural-language processing research can be stated quite simply: to endow computers with human language capability. The pursuit of this objective, however, has been a difficult task for at least two reasons: first, this capability is far from being a well-understood phenomenon; second, the tools for teaching computers what we do know about human language are still very primitive. The solution of these problems lies within the respective domains of linguistics and computer science.

Similar problems have arisen in computer science. Whenever a new computer application area emerges, there follow new modes of communication with computers that are geared towards such areas. Computer languages are a direct result of this need for effective communication with computers. A considerable body of accumulated knowledge about the design of languages for communicating information to computers has been derived from the subfields of programming language design and semantics.

2. The Critical Properties of the Language

It is not the purpose of this paper to provide a comprehensive description of the PATR-II project, or even of the formalism itself. Rather, we will discuss briefly the critical properties. It has been the goal of the PATR group at SRI to utilize a relevant portion of this knowledge in implementing tools to facilitate communication of linguistic information to computers.

The PATR-II formalism is our current computer language for encoding linguistic information. This paper, a brief overview of that formalism, attempts to explicate our design decisions in terms of a set of properties that effective computer languages should incorporate, namely: simplicity, power, mathematical well-foundedness, flexibility, implementability, modularity, and declarativeness. More extensive discussions of various aspects of the PATR-II formalism and systems can be found in papers by Shieber et al., [83], Pereira and Shieber [84] and Karttunen [84].

The notion of designing specialized computer languages and systems to encode linguistic information is not new; PROGRAMMAR [Winograd, 72], ANNs [Woods, 70], and DIALOGIC [Grosz, et al., 82] are but a few of the better-known examples. Furthermore, a trend has arisen recently in linguistics towards declarativeness in grammar formalisms—for instance, lexical-functional grammar (LFG) [Bresnan, 83], generalized phrase-structure grammar (GPSG) [Gazdar and Pullum, 82] and functional unification grammar (UG) [Kay, 83]. Finally, in computer science there has been a great deal of interest in declarative languages (e.g., logic programming and specification languages), and their supporting denotational semantics. But to our knowledge, no attempt has yet been made to combine the three approaches so as to yield a declarative computer language with clear semantics designed specifically for encoding linguistic information. Such a language, of which PATR-II is an example, would reflect a felicitous convergence of ideas from linguistics, artificial intelligence, and computer science.
properties of PATR-II to give a flavor for our approach to
the design of the language. References to papers with more
complete descriptions of particular aspects of the project
are provided when appropriate.

2.1. Simplicity: An Introduction to the
PATR-II Formalism

Building on a convergence of ideas from the linguistics
and AI communities, PATR-II takes as its primitive opera-
tion an extended pattern-matching technique, unification,
first used in logic and theorem-proving research and lately
finding its way into research in linguistics [Kay, 79; Gazdar
and Pullum, 82] and knowledge representation [Reynolds,
70; Ait-Kaci, 83]. Instead of unifying logic terms, how-
ever, PATR unification operates on directed acyclic graphs
(DAGs).

DAGs can be atomic symbols or sets of label/value
pairs, where the values are themselves DAGs (either atomic
or complex). Two labels can have the same value—thus the
use of the term graph rather than tree. DAGs are notated
either by drawing the graph structure itself, with the la-
bels marking the arcs, or, as in this paper, by notating the
sets of label/value pairs in square brackets, with the labels
separated from their values by a colon; e.g., a DAG associ-
ated with the verb “knight” (as in “Uther wants to knight
Arthur”) would appear (in at least one of our grammars) as

\[
\text{[cat: v} \\
\text{head: [aux: false}} \\
\text{form: nonfinite} \\
\text{voice: active} \\
\text{trans: [pred: knight}} \\
\text{arg1: [f1134]}} \\
\text{arg2: [f1138]}} \\
\text{syncat: [first: [cat: np}} \\
\text{head: [trans: [f1134]]}} \\
\text{rest: [first: [cat: np}} \\
\text{head: [trans: [f1138]]}} \\
\text{rest: [f1160]}} \\
\text{tail: [f1160]]}
\]

Reentrant structure is notated by labeling the DAG with
an arbitrary label (in angle brackets), then using that label
for future references to the DAG.

Associated with each entry in the lexicon is a set of
DAGs.* The root of each DAG will have an arc labeled cat
\(^{3}\) Technically, these are rooted, directed, acyclic graphs with labeled
arcs. Formal definition of these and other technical notions can be
found in Appendix A of Shieber et al. [83]. Note that some imple-
mentations have been extended to handle cyclic graph structures as
well as graph structures with disjunction and negation [Karttunen,
84].

\(^{4}\) In our implementation, this association is not directly encoded—since
this would yield a grossly inefficient characterization of the lexicon—
but is mediated by a morphological analyzer. See Section 2.6 for
further details.

whose value will be the category of the associated lexical
entry. Other arcs may encode information about the synt-
tactic features, translation, or syntactic subcategorization
of the entry. But only the label cat has any special sig-
ificance; it provides the link between context-free phrase
structure rules and the DAGs, as explicated below.

PATR-II grammars consist of rules with a context-free
phrase structure portion and a set of unifications on the
DAGs associated with the constituents that participate in
the application of the rule. The grammar rules describe how
constituents can be built up to form new constituents with
associated DAGs. The right side of the rule lists the cat
values of the DAGs associated with the filial constituents;
the left side, the cat of the parent. The associated uni-
fications specify equivalences that must exist among the
various DAGs and sub-DAGs of the parent and children.
Thus, the formalism uses only one representation—DAGs—
for lexical, syntactic, and semantic information, and one
operation—unification—on this representation.

By way of example, we present a trivial grammar for a
fragment of English with a lexicon associating words with
DAGs.

\[
S \rightarrow NP \ VP \\
\text{\textless VP agr\textgreater} = \text{\textless NP agr\textgreater} \\
\text{\textless VP agr\textgreater} = \text{\textless V agr\textgreater} \\
\text{Uther:} \\
\text{\textless cat\textgreater} = \text{np} \\
\text{\textless agr number\textgreater} = \text{singular} \\
\text{\textless agr person\textgreater} = \text{third} \\
\text{Arthur:} \\
\text{\textless cat\textgreater} = \text{np} \\
\text{\textless agr number\textgreater} = \text{singular} \\
\text{\textless agr person\textgreater} = \text{third} \\
\text{knight:} \\
\text{\textless cat\textgreater} = \text{v} \\
\text{\textless agr number\textgreater} = \text{singular} \\
\text{\textless agr person\textgreater} = \text{third}
\]

This grammar (plus lexicon) admits the two sentences
“Uther knights Arthur” and “Arthur knights Uther.” The
phrase structure associated with the first of these is:

\[
\text{[s [NP Uther] [VP [v knight] [NP Arthur]]]}
\]

The VP rule requires that the agr feature of the DAG
associated with the VP be the same as (unified with) the agr
of the V. Thus, the VP’s agr feature will have as its value
the same node as the V’s agr, and hence the same values for
the person and number features. Similarly, by virtue of the
unification associated with the S rule, the NP will have
the same agr value as the VP and, consequently, the V. We
have thus encoded a form of subject-verb agreement.

Note that the process of unification is order-independent.
For instance, we would get the same effect regardless of
whether the unifications at the top of the parse tree were
effected before or after those at the bottom. In either case,
the DAG associated with, e.g., the VP node would be
As a simple example, we could define (with the following syntax) the template Verb as

Let Verb be

\[<\text{cat}> = V\]

and the template ThirdSing as

Let ThirdSing be

\[<\text{agr number}> = \text{singular}\]
\[<\text{agr person}> = \text{third}\]

The lexical entry for “knights” would then be temporarily sacrificed in an attempt to keep the underlying formalism simple, general, and semantically well-founded. However, given a simple underlying formalism, we can build more efficient, specialized languages on top of it, much as MACLISP might be built on top of pure LISP. And just as MACLISP need not be implemented (and is not implemented) directly in pure LISP, specialized formalisms built conceptually on top of pure PATR-II need not be so implemented (although currently we do implement them directly through pure PATR-II). The effectiveness of this approach can be seen in the fact that at least a sizable portion of English syntax has been encoded in various experimental PATR-II grammars constructed to date. The syntactic constructs encoded include subcategorization of various complement types (NPs, Ss, etc.), active, passive, “there” insertion, extraposition, raising, and equi-NP constructions, and unbounded dependencies (such as Wh-movement and relative clauses). Other theory-dependent devices that have been modeled with PATR-II include head-feature percolation [Gazdar and Pullum, 82], and LFG-like semantic forms [Kaplan and Bresnan, 83]. Note that none of these constructs and techniques required expansion of the underlying formalism; indeed, the constructions all make use of the techniques described in this section. See Shieber et al. [83] for a detailed discussion of the modeling of some of these phenomena.

The devices now available for molding PATR-II to conform to a particular intended usage or linguistic theory are in their nascent stage. However, because of their great importance in making the PATR-II system a usable one, we will discuss them briefly. It is important to keep in mind that these methods should not be considered a part of the underlying formalism, but merely “syntactic sugar” to increase the system's utility and allow it to conform to a user's intentions.

2.4.1. Templates

Because so much of the information in the PATR-II grammars under actual development tends to be encoded in the lexicon, most of our research has been devoted to methods for removing redundancy in the lexicon by allowing the users themselves to define primitive constructs and operations on lexical items. Primitive constructs, such as the transitive, dyadic, or equi-NP properties of a verb, can be defined by means of templates, that is, DAGs that encode some linguistically isolable portion of the DAG of a lexical item. These template DAGs can then be combined to build the lexical item out of the user-defined primitives.

As a simple example, we could define (with the following syntax) the template Verb as

Let Verb be

\[<\text{cat}> = V\]

and the template ThirdSing as

Let ThirdSing be

\[<\text{agr number}> = \text{singular}\]
\[<\text{agr person}> = \text{third}\]

The lexical entry for “knights” would then be

---

\(\text{[cat: vp} \quad \text{agr: [person: third}}\)
\(\text{number: singular]}}\)
Templates can themselves refer to other templates, enabling definition of abstract linguistic concepts hierarchically. For instance, a modal verb template may use an auxiliary verb template, which in turn may be defined using the verb template above. In fact, templates are currently employed for abstracting notions of subcategorization, verb form, semantic type, and a host of other concepts.

2.4.2. Lexical Rules

More complex relationships among lexical items can be encoded by means of lexical rules. These rules, such as passive and "there" insertion, are user-definable operations on the lexical items, enabling one variant of a word to be built from the specification of another variant. A lexical rule is specified as a set of selective unifications relating an input DAG and an output DAG. Thus, unification is the primitive used in this device as well.

Lexical rules are used to encode the relationships among various lexical entries that would typically be thought of as transformations or relation-changing rules (depending on one's ideological outlook). Because lexical rules perform these operations, the lexicon need include only a prototype entry for each verb. The variant forms can be derived through lexical rules applied in accordance with the morphology actually found on the verb. (The morphological analysis in the implementations of PATR-II is performed by a program based on the system of Koskenniemi [83] and was written by Lauri Karttunen [83].)

For instance, given a PATR-II grammar in which the DAGs are used to emulate the f-structures of LFG, we might write a passive lexical rule as follows (following Brennan [83]):

Define Passive as

\[
\begin{align*}
\text{<out cat>} &= \text{<in cat>} \\
\text{<out form>} &= \text{passprt} \\
\text{<out subj>} &= \text{<in obj>} \\
\text{<out obj>} &= \text{<in subj>}
\end{align*}
\]

The rule states in effect that the output DAG (the one associated with the passive verb form) marks the lexical item as being a passive verb whose object is the input DAG's subject and whose subject is the input's object. Such lexical rules have been used for encoding the active/passive dichotomy, "there" insertion, extraposition, and other so-called relation-changing rules.

2.5. Modularity and Declarativeness

The PATR-II formalism is a completely declarative formalism, as evidenced by its denotational semantics and the order-independence of its definition. Modularity is achieved through the ability to define primitive templates and lexical rules that are shared among lexical items, as well as by the declarative nature of the grammar formalism itself,

removing problems of interaction of rules. Rules are guaranteed to always mean the same thing, regardless of the environment of other rules in which they are placed.

2.6. Implementability

Implementability is an empirical matter, given credence by the fact that we now have three implementations of the formalism. One desirable aspect of the simplicity and declarative nature of the formalism is that even though the three implementations differ substantially from one another, using different parsing algorithms (with both top down and bottom up properties), different implementations of unification, different methods of compiling the rules, all are able to run on exactly the same grammars yielding the identical results.

The three implementations of the PATR-II system currently in operation at SRI are as follows:

- An INTERLISP version for the DEC-2060 using a variant of the Cocke-Kasami-Younger parsing algorithm and the KIMMO morphological analyzer [Karttunen, 83], and a limited programming environment.
- A ZETALISP version for the Symbolics 3600 using a left-corner parsing algorithm and the KIMMO morphological analyzer, with an extensive programming environment (due primarily to Mabry Tyson) that includes incremental compilation, multiple window debugging facilities, tracing, and an integrated editor.
- A Prolog version (DEC-10 Prolog) running on the DEC-2060 by Fernando Pereira, designed primarily as a testbed for experimentation with efficient structure-sharing DAG unification algorithms, and incorporating an Earley-style parsing algorithm.

In addition, Lauri Karttunen and his students at the University of Texas have implemented a system based on PATR-II but with several interesting extensions, including disjunction and negation in the graph structures [Karttunen, 84]. These extensions will undoubtedly be integrated into the SRI systems and formal semantics for them are being pursued.

3. Conclusion

The PATR-II formalism was designed as a computer language for encoding linguistic information. The design was influenced by current theory and practice in computer science, and especially in the areas of programming language design and semantics. The formalism is simple (consisting of just one primitive operation, unification), powerful (although it can be constrained to be decidable), mathematically well-founded (with a complete denotational semantics), flexible (as demonstrated by its ability to model analyses in GPSG, LFG, DCG and other formalisms), modular (because of its higher-level notational devices such as templates and lexical rules), declarative (yielding order-independence of operations), and implementable (as demonstrated by three quite dissimilar implemented systems and one highly developed programming environment).
As we have emphasized herein, PATR-II seems to represent a convergence of techniques from several domains—computer science, programming language design, natural language processing and linguistics. Its positioning at the center of these trends arises, however, not from the admixture of many discrete techniques, but rather from the application of a single simple yet powerful concept to the encoding of linguistic information.

References


