Evolutionary Games on Cycles

Citation

Published Version
doi:10.1098/rspb.2006.3576

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Accessibility
Evolutionary games on cycles

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Traditional evolutionary game theory explores frequency dependent selection in well-mixed populations without spatial or stochastic effects. But recently there has been much interest in studying evolutionary game dynamics in spatial settings, on lattices and other graphs. Here we present an analytic approach for stochastic evolutionary game dynamics on the simplest possible graph, the cycle. For three different update rules, called ‘birth-death’, ‘death-birth’ and ‘imitation’, we derive exact conditions for natural selection to favor one strategy over another. As specific examples, we consider a coordination game and the Prisoner’s Dilemma. In the latter, selection can favor cooperators over defectors for ‘death-birth’ updating and ‘imitation’. We also study the case where the replacement graph of evolutionary updating remains a cycle, but the interaction graph for playing the game is a complete graph. In this setting, all three update rules lead to identical conditions in the limit of weak selection, where we find the ‘1/3-law’ of well mixed populations.

Keywords: evolutionary dynamics, social networks, fixation probability, spatial games, evolutionary graph theory, prisoner’s dilemma, frequency dependent selection

1. INTRODUCTION

Game theory was invented by von Neumann and Morgenstern (1944) as a mathematical approach to understand strategic and economic decisions of humans. Hamilton (1967),
Trivers (1971) and Maynard Smith & Price (1973) brought game theory to biology. Instead of analyzing the interaction between two rational players, evolutionary game theory explores the dynamics of a population of players under the influence of natural selection (Maynard Smith 1982). In the classical setting of the replicator equation, the population size is infinite and interactions are equally likely between any two individuals (Taylor & Jonker 1978, Hofbauer et al 1979, Zeeman 1980). Each individual obtains an average payoff which is interpreted as biological fitness: strategies reproduce according to their payoff. Successful strategies spread and eliminate less successful ones. The payoff depends on the frequency of strategies in the population. Hence, natural selection is frequency dependent. The replicator equation is deeply connected to the concept of an evolutionarily stable strategy (ESS) or Nash equilibrium. In the framework of the replicator equation, an ESS cannot be invaded by any mutant strategy (Hofbauer & Sigmund 1998). For recent books on game theory and evolutionary game theory we refer to Fudenberg & Tirole (1991), Binmore (1994), Weibull (1995), Samuelson (1997), Fudenberg & Levine (1998), Hofbauer & Sigmund (1998), Gintis (2000), and Cressman (2003). Recent reviews of evolutionary game dynamics are Hofbauer & Sigmund (2003) and Nowak & Sigmund (2004).


Ellison (1993) has studied coordination games in spatial settings and found that
localized interactions facilitate faster convergence than global interactions. Nakamaru et al (1997, 1998) have studied the interaction between tit-for-tat and always-defect in the repeated Prisoner’s Dilemma on lattices. They have observed that ‘fertility’ selection is less favorable for cooperation than ‘mortality’ selection. Nakamaru and Iwasa (2005) analyze the evolution of altruism in a spatially structured population with punishment (Fehr & Gaechter 2000, Sigmund et al 2001).


In evolutionary games in finite populations, stochasticity plays an important role (Nowak et al 2004, Taylor et al 2004, Traulsen et al 2005, Imhof & Nowak 2006). The most crucial quantities are the fixation probabilities. If a single individual playing a strategy $A$ is added to a population playing strategy $B$, then there is a certain probability that $A$ will generate an offspring lineage which will eventually take over the entire population. For a neutral mutant, this fixation probability is given by the inverse of the population size, $1/N$. If the fixation probability of $A$ is greater than $1/N$, then selection favors the fixation of this strategy. The analysis of evolutionary game dynamics in finite populations requires a comparison of the fixation probabilities of the two strategies with each other and with $1/N$.

Lieberman et al (2005) show how to calculate the fixation probability of a randomly placed mutant on various graphs. The traditional well-mixed population is the special case of a complete graph (all individuals are connected) with identical weights. For constant (instead of frequency dependent) selection, all graphs which are circulations lead to the same fixation probabilities as the complete graph. A graph is a circulation if for each vertex the sum over all incoming weights equals the sum over all outgoing weights. Lieberman et al (2005) also discuss frequency dependent selection (evolutionary games) on directed cycles.
Ohtsuki et al (2006) study the evolution of cooperation on a large variety of graphs. Using pair approximation (Matsuda et al 1987, 1992, Harada et al 1995), they find that selection favors the evolution of cooperation if \( b/c > k \). The benefit to cost ratio of the altruistic act has to exceed the (average) number of neighbors per individual. Computer simulations suggest that this very simple rule works well for many different graphs including lattices, regular graphs, random regular graphs and scale-free networks. Of course, pair approximation makes certain assumptions that may or may not hold in a particular setting. Therefore, the attempt of this paper is to derive exact results for the fixation probabilities of evolutionary games on a simple family of graphs, the cycle. In these cases, direct calculations succeed, because a single invader always leads to one cluster that does not fragment into pieces. A cycle is a regular graph with \( k = 2 \); each individual has exactly two neighbors. Therefore, we expect that \( b/c > 2 \) will be a decisive rule for the evolution of cooperation on cycles.

In Section 2, we introduce the basic rules of the game, consider three different update mechanism for the evolutionary dynamics and present results for general two person games. In Section 3, we study coordination games on cycles. In Section 4, we study the evolution of cooperation. Section 5 assumes that the updating (or the evolutionary competition) still occurs between nearest neighbor on a cycle, but the payoff of each individual is derived from random interactions among the whole population. In this case, we find that all three update mechanisms lead to the same criterion. Throughout the paper we state the crucial results in the main text and show their derivations in the Appendix.

### 2. GAMES ON CYCLES

Consider a game between two strategies \( A \) and \( B \). Denote by \( a, b, c \) and \( d \) the payoff for \( A \) versus \( A \), \( A \) versus \( B \), \( B \) versus \( A \) and \( B \) versus \( B \), respectively. The payoff matrix is given by

\[
\begin{pmatrix}
A & B \\
A \begin{pmatrix}
a & b \\
c & d
\end{pmatrix}
\end{pmatrix}
\]

(1)

Strategy \( A \) is a strict Nash equilibrium if \( a > c \). In this case, strategy \( A \) is also evolutionarily stable, which means that \( B \) cannot invade \( A \) in an infinitely large, well-mixed
population. Similarly, strategy $B$ is a strict Nash equilibrium and evolutionarily stable if $b < d$. If $a > c$ and $b < d$ then both strategies are strict Nash equilibria. In this case, a well-mixed population will either converge to one strategy or the other depending on the initial frequencies of the strategies. The strategy with the bigger basin of attraction is called ‘risk-dominant’. A simple calculation shows that $A$ is risk dominant if $a + b > c + d$.

If $a > c$ and $b > d$ then $A$ dominates $B$, which means that any mixed population will converge to a homogenous state using only strategy $A$. All these results are based on deterministic evolutionary game dynamics in infinitely large, well-mixed populations without spatial effects.

Real biological populations are neither well-mixed nor infinite. Let us therefore consider a structured, finite population with individuals $i = 1, \ldots, N$. The rate at which individual $i$ plays the game with $j$ is given by $h_{ij}$. The rate at which an offspring of individual $i$ replaces $j$ is given by $g_{ij}$. The matrix $G = [g_{ij}]$ specifies the ‘replacement graph’, while the matrix $H = [h_{ij}]$ specifies the ‘interaction graph’ (Fig. 1). In a series of other papers (Ohtsuki et al 2006, Ohtsuki & Nowak 2006), we have used pair-approximation and computer simulation to study evolutionary games on a large variety of graphs. Here instead we derive exact results for the simple case that $G$ and $H$ are identical and given by cycles with equal weights.

Each player is interacting with its two immediate neighbors. The payoffs from these two interactions are added up. A parameter $w$ measures the intensity of selection. The fitness of an individual is given by $1 - w + wP$, where $P$ is the individual’s payoff. The case $w = 1$ denotes strong selection: fitness equals payoff. The case $w \ll 1$ denotes weak selection: the payoff from the game represents only a small contribution to fitness. For $w = 0$, we obtain neutral drift; all strategies have the same fitness. It is interesting to note that for the traditional replicator equation, the intensity of selection cancels out, but for stochastic game dynamics the intensity of selection plays an important role (Nowak et al 2004). Often simple and illuminating results arise in the limit of weak selection, $w \to 0$.

We consider three different update rules: (i) ‘birth-death’ (BD) means that an individual is selected for reproduction proportional to fitness and the offspring replaces a randomly chosen neighbor; (ii) ‘death-birth’ (DB) means that a random individual is
eliminated, and the neighbors compete for the empty site proportional to their fitness; (ii) ‘imitation’ (IM) means that a random individual is chosen to update its strategy; it will either stay with its own strategy or imitate one of the neighbors’ strategy proportional to fitness.

Imagine now that a single A individual is added to a population of N − 1 B individuals. The A individual could die before reproducing or produce a lineage of A, which becomes extinct after some time. In both cases, the population returns to a state of ‘all-B’. The other possibility is that A produces a lineage which will eventually take over the entire population, which means that B becomes extinct. In this case, the population will end up in the state ‘all-A’. Denote by ρ_A the probability that a single A individual will take over a population of B. Denote by ρ_B the probability that a single B individual will take over a population of A. The quantities ρ_A and ρ_B are called the fixation probabilities of A and B, respectively.

A neutral mutant has fixation probability 1/N. Therefore, if ρ_A > 1/N then selection favors the fixation of A. If ρ_A < 1/N then selection opposes the fixation of A. If ρ_A > ρ_B then selection favors A over B.

For strong selection (w = 1) and large population size (N → ∞), we find that selection favors the fixation of A and opposes the fixation of B, ρ_A > 1/N > ρ_B, if

\[
\begin{align*}
\text{BD} & : \quad a + b > c + d \\
\text{DB} & : \quad a(a + b) > (c + d)d \\
\text{IM} & : \quad (3a + b)(a + b) > (c + d)(c + 3d)
\end{align*}
\]

(2)

For the limit of weak selection (w → 0) and large population size (N → ∞), we find that ρ_A > 1/N > ρ_B, if

\[
\begin{align*}
\text{BD} & : \quad a + b > c + d \\
\text{DB} & : \quad 3a + b > c + 3d \\
\text{IM} & : \quad 5a + 3b > 3c + 5d
\end{align*}
\]

(3)

Interestingly, for BD updating, we find the same condition for weak and strong selection, and this condition is equivalent to risk-dominance. For DB and IM updating,
however, there are different conditions for weak and strong selection, and neither of those
conditions are equivalent to risk dominance.

In the limit of weak selection, we can also derive simple conditions for any given
population size, \( N \). We find that \( \rho_A > \rho_B \) if

\[
\begin{align*}
\text{BD} & : (N - 2)a + Nb > Nc + (N - 2)d \\
\text{DB} & : (3N - 8)a + Nb > Nc + (3N - 8)d \\
\text{IM} & : (5N - 12)a + 3Nb > 3Nc + (5N - 12)d
\end{align*}
\]

For large \( N \), we recover the conditions given by (3). In contrast to (2) and (3), however,
inequalities (4) do not specify whether \( \rho_A \) or \( \rho_B \) are greater or smaller than \( 1/N \). These
conditions are derived in the Appendix.

3. PARETO EFFICIENCY AND RISK DOMINANCE

Consider a coordination game. Strategies \( A \) and \( B \) are strict Nash equilibria. We
have \( a > c \) and \( b < d \). Suppose \( B \) is risk-dominant, which means that \( a + b < c + d \). But
\( A \) is Pareto-efficient, which means that \( a > d \). Hence, strategy \( B \) has the bigger basin of
attraction, but a homogeneous population of \( A \) has a higher payoff than a homogeneous
population of \( B \).

For BD updating, we always have that \( \rho_A < \rho_B \), which means that the fixation
probability of \( A \) is smaller than that of \( B \). BD updating on the cycle favors risk-dominance
over Pareto-efficiency. For DB and IM updating, however, it is possible that the Pareto
efficient strategy is favored by selection. A particular example is given by the payoff matrix

\[
\begin{pmatrix}
A & B \\
A & 10 \\
B & 9
\end{pmatrix}
\] (5)

Here \( A \) is pareto-efficient, while \( B \) is risk-dominant. Nevertheless, the conditions (2) and
(3) hold for DB and IM updating, which implies that \( \rho_A > 1/N > \rho_B \) for large \( N \).
Therefore, these two update rules can lead to efficient outcomes for coordination games on
cycles.

4. THE PRISONER’S DILEMMA ON A CYCLE
As a specific example, consider the Prisoner’s Dilemma, which is given by the payoff matrix

\[
\begin{pmatrix}
C & D \\
C & R & S \\
D & T & P
\end{pmatrix}
\]

(6)

The ranking of the payoff values is \( T > R > P > S \). The ‘temptation’ to defect, \( T \), is greater than the ‘reward’ for mutual cooperation, \( R \), which is greater than the ‘punishment’, \( P \), for mutual defection, which is greater than the ‘sucker’s payoff’, \( S \). Since \( T > R \) and \( P > S \), cooperators, \( C \), are dominated by defectors, \( D \). The traditional replicator dynamics and stochastic dynamics in a well-mixed population of finite size (Nowak et al, 2004) both lead to defection. Moreover, we have \( R + S < T + P \), which means that BD updating on the cycle also favors defectors over cooperators.

But for DB or IM updating, it is possible that cooperators are favored over defectors. A numerical example is given by

\[
\begin{pmatrix}
C & D \\
C & 10 & 1 \\
D & 11 & 2
\end{pmatrix}
\]

(7)

For this payoff matrix, the fixation probability of a single cooperator, \( \rho_C \), is greater than \( 1/N \), while the fixation probability of a single defector, \( \rho_D \), is less than \( 1/N \), for strong and weak selection, for DB and IM updating. Therefore, evolutionary game dynamics on cycles can favor cooperation over defection in the Prisoner’s Dilemma.

The intuitive reason for this finding is the following. On a cycle an invasion attempt by a single cooperator leads to a single cluster of cooperators. The question is whether the boundary between cooperators and defectors moves in favor of cooperators or not. For BD updating only the payoff of the two individuals (one cooperator and one defector) right at the boundary affect the stochastic dynamics. The cooperator always has a lower payoff than the defector. For DB and IM updating, the payoffs of the two individuals that are one place removed from the boundary also affect the dynamics. The ‘cooperator once removed’ always has a higher payoff than the ‘defector once removed’. Therefore, it is possible that the boundary moves in favor of cooperators (Fig 2).
For a simplified Prisoner’s Dilemma given by two parameters denoting the cost, $c$, and benefit, $b$, of an altruistic act, we obtain the payoff matrix

$$
\begin{pmatrix}
C & D \\
C & (b - c) \\
D & b \\
\end{pmatrix}
$$

(8)

The conditions for strong selection (eq 2) cannot be used for this payoff matrix, because our analysis requires strictly positive fitness values. The conditions for weak selection (eq 3) can be used, and here we obtain particularly elegant results. For BD updating we find that $\rho_D > 1/N > \rho_C$ for any choice of $b, c$ and $N$. Hence, for BD updating, cooperation cannot evolve. For the other update mechanisms, however, we find that $\rho_C > 1/N > \rho_D$ if

$$
\text{DB : } \frac{b}{c} > 2 + \frac{4}{N - 4} \\
\text{IM : } \frac{b}{c} > 4 + \frac{18}{N - 6}
$$

(9)

Note that the DB condition requires $N > 4$, while IM requires $N > 6$. For large $N$, we obtain

$$
\text{DB : } \frac{b}{c} > 2 \\
\text{IM : } \frac{b}{c} > 4
$$

(10)

These findings concur with the results obtained by pair-approximation for regular graphs of degree $k$ (Ohtsuki et al 2006). There we find that DB updating favors cooperators if $b/c > k$, while IM updating favors cooperators if $b/c > k + 2$. Both results hold for large $N$ and weak selection. In this framework, we also find that BD updating never favors cooperators. Since the cycle has degree $k = 2$, the findings are in agreement.

5. LOCAL UPDATING BUT GLOBAL INTERACTION

Let us return to the general game between two strategies given by payoff matrix (1). Let us keep the cycle as the replacement graph, but use a complete interaction graph. This means that updating occurs locally among nearest neighbors, but interactions are globally. Remarkably, the following results hold for all three update mechanisms.
In the limit of weak selection, we find that \( \rho_A > \rho_B \) if

\[
(N - 2)a + Nb > Nc + (N - 2)d.
\]  

(11)

This inequality leads to \( a + b > c + d \) for large \( N \). Furthermore, for large \( N \), we find that \( \rho_A > 1/N \) if

\[
a + 2b > c + 2d.
\]  

(12)

This condition is the ‘1/3-rule’ which has been observed previously for well-mixed populations. If strategies \( A \) and \( B \) are best replies to themselves, then the standard replicator equation has an unstable equilibrium at a frequency of \( A \) given by \( x^*_A = (d-b)/(a-b-c+d) \).

For stochastic game dynamics in a well-mixed population using the Moran process, the fixation probability \( \rho_A \) exceeds \( 1/N \) for weak selection and large population size if \( x^*_A < 1/3 \) (Nowak et al 2004). This ‘1/3-rule’ has also been obtained for frequency dependent selection in the Wright-Fisher process (Imhof & Nowak 2006).

6. CONCLUSION

We have studied stochastic evolutionary game dynamics on cycles, which represent the simplest possible family of graphs. Each individual has two neighbors. A single invader leads to a single cluster, which does not fragment into pieces. Therefore, the transition matrix of the stochastic process is tri-diagonal, and the fixation probability of an invading strategy can be calculated explicitly. We consider three different update mechanisms for evolutionary dynamics, which we call ‘birth-death’ (BD), ‘death-birth’ (DB) and ‘imitation’ (IM). We explore the interaction between two strategies \( A \) and \( B \) and present conditions for the fixation probability \( \rho_A \) to be greater than \( \rho_B \). We also compare the two fixation probabilities with \( 1/N \), which is the fixation probability of a neutral mutant. For strong and weak selection, we find simple conditions (eqs 2 and 3) for large population size, \( N \). For weak selection, we also find simple conditions for any population size (eq 4).

Note that none of these conditions reduce to the simple criteria of a strict Nash equilibrium or an evolutionarily stable strategy (ESS). Therefore, if \( A \) is a strict Nash
equilibrium or an ESS, then it can still be the case that $\rho_A$ is less than $\rho_B$ and/or less than $1/N$.

In coordination games, BD updating always favors the risk dominant strategy, while DB and IM updating can favor Pareto efficiency over risk dominance. In the Prisoner’s Dilemma, BD updating always favors defectors, while DB and IM updating can favor cooperators over defectors. In the latter case, a cooperator lineage starting from a single cooperator invades a population of defectors with a probability of an advantageous mutant. If cooperators pay a cost $c$ for neighbors to receive a benefit $b$, then DB updating favors the evolution of cooperation if $b/c > 2$. For IM updating, we find $b/c > 4$. These findings are in agreement with the results obtained by pair approximation for regular graphs (Ohtsuki et al 2006).

Note that these conditions are needed for the evolution of cooperation by ‘spatial reciprocity’ alone (Nowak & May 1992) without any strategic complexity (Axelrod & Hamilton 1981, Axelrod 1984) or reputation effects (Nowak & Sigmund 2005). One can also study the synergistic interaction between spatial reciprocity and direct or indirect reciprocity, which will lead to less stringent conditions for the emergence and stability of cooperation.

The intuitive reason for the different behavior of the update mechanisms is illustrated in Figure 2. For BD updating, only the payoff of the two individuals right at the edge of a cluster plays a role, because two adjacent players directly compete for reproduction. The scale of interaction and competition is the same there. In contrast, for DB and IM updating two individuals that are two-steps away from the edge of the cluster are also involved in the competition; therefore, the scale of competition is larger than that of interaction. For the struggle between cooperators and defectors, this is a crucial difference, because right on the edge, the defector always has a higher payoff than the cooperator, but the cooperator that is once removed from the edge can have a higher payoff still. There is an interesting parallel between this observation and previous studies on viscous populations by Taylor 1992 and West et al 2002, who show that cooperation is favored when the scale of competition is larger than the scale of interaction.

Studying games on graphs leads to an understanding of how population structure af-
ffects the outcome of evolutionary dynamics. Compared to most other evolutionary graphs, the cycle has the particular advantage that the fixation probabilities in frequency depend-ent selection can be calculated exactly. Moreover, the cycle represents the extreme case where the effect of spatial structure on evolutionary dynamics is strongest. On the other end of the spectrum is the well-mixed population, which is given by a complete graph (where all individuals are connected to each other). For the Prisoner’s Dilemma, this means that the condition $b/c > 2$ is a minimum requirement. All other structures will demand a larger benefit-to-cost ratio for the evolution of cooperation.

Acknowledgments

Support from the John Templeton Foundation, the Japan Society for the Promotion of Science and KAKENHI is gratefully acknowledged. The Program for Evolutionary Dynamics at Harvard University is sponsored by Jeffrey Epstein.

Appendix A. DERIVATION OF THE RESULTS IN THE MAIN TEXT

Let us calculate $\rho_A$.

If the replacement graph is a cycle, then a single invader always leads a single connected cluster. Therefore, in order to describe the stochastic process, it is sufficient to count the number of players. Let $i$ denote the number of $A$ players. The state space of the stochastic process is given by $i = 0, ..., N$, where $N$ is the total population size, which is constant. At each time step, the variable $i$ can at most change by one unit. Denote the associated transition probabilities by $\lambda_i = \text{Prob.}(i \to i + 1)$ and $\mu_i = \text{Prob.}(i \to i - 1)$. We have $1 - \lambda_i - \mu_i = \text{Prob.}(i \to i)$.

There are two absorbing states $i = N$ and $i = 0$. The former corresponds to ‘all-$A$’ and the latter to ‘all-$B$’. If the system has reached such a state it will stay there forever. All other states are transient. Denote by $\rho_A$ the probability to reach state $i = N$ when starting from $i = 1$. This quantity is the fixation probability of $A$. Similarly, the fixation probability of $B$, $\rho_B$, denotes the probability to reach state $i = 0$ when starting from state $i = N - 1$. A direct calculation (Karlin & Taylor 1975, Ewens 2004) shows that

$$\rho_A = 1 \left/ \left(1 + \sum_{j=1}^{N-1} \prod_{i=1}^{j} \frac{\mu_i}{\lambda_i} \right) \right.. \quad (A.1)$$
For the ratio of the two fixation probabilities, we find
\[
\frac{\rho_A}{\rho_B} = \prod_{i=1}^{N-1} \frac{\lambda_i}{\mu_i}. \tag{A.2}
\]
Thus we need to calculate \(\mu_i/\lambda_i\) for each \(i\).

For convenience, define \(\alpha, \beta, \gamma\) and \(\delta\) as
\[
\alpha = 1 - w + 2wa, \beta = 1 - w + 2wb, \gamma = 1 - w + 2wc \quad \text{and} \quad \delta = 1 - w + 2wd.
\]
Here \(a, b, c, d\) are the entries of the payoff matrix and \(w\) denotes the intensity of selection.

(i) For BD updating we find
\[
\frac{\mu_i}{\lambda_i} = \left\{ \begin{array}{ll}
\frac{(\gamma + \delta) / 2\beta}{\mu_i} & (i = 1) \\
\frac{(\gamma + \delta) / (\alpha + \beta)}{2\gamma / (\alpha + \beta)} & (2 \leq i \leq N - 2) \\
\frac{(\gamma + \delta) / (\alpha + \beta)}{2\gamma / (\alpha + \gamma)} & (i = N - 1)
\end{array} \right.
\]
For strong selection \((w = 1)\), the relative size of \((\gamma + \delta) / (\alpha + \beta)\) compared to one is important. If it is less than one, which is the case when \(a + b > c + d\), then \(\rho_A\) is positive, otherwise \(\rho_A = 0\). For weak selection, we have
\[
\rho_A = \frac{1}{N} + \frac{w}{2N^2} \left\{ (N^2 - 3N + 2)a + (N^2 + N - 2)b - (N^2 - N + 2)c - (N^2 - N - 2)d \right\}
\]
and
\[
\frac{\rho_A}{\rho_B} = 1 + w \{ (N - 2)a + Nb - Nc - (N - 2)d \}.
\]

(ii) For DB updating we obtain
\[
\frac{\mu_i}{\lambda_i} = \left\{ \begin{array}{ll}
\frac{(\beta + \delta) / 2\beta}{\mu_i} & (i = 1) \\
\frac{[(\gamma + \delta) / (\alpha + \beta)] \cdot [(\alpha + \beta + 2\delta) / (\alpha + \beta + \gamma + \delta)]}{[2\gamma / (\alpha + \gamma)]} & (2 \leq i \leq N - 3) \\
\frac{[(\gamma + \delta) / (\alpha + \beta)] \cdot [(\alpha + \beta + \gamma + \delta) / (2\alpha + \gamma + \delta)]}{[2\gamma / (\alpha + \gamma)]} & (i = N - 2) \\
\frac{[(\gamma + \delta) / (\alpha + \beta)]}{[2\gamma / (\alpha + \gamma)]} & (i = N - 1)
\end{array} \right.
\]
For strong selection, \(\rho_A > 0\) if and only if \(\frac{[(\gamma + \delta) / (\alpha + \beta)] \cdot [(\alpha + \beta + 2\delta) / (2\alpha + \gamma + \delta)]}{13} < 1\), which is equivalent to \(a(a + b) > (c + d)d\). For weak selection, we have
\[
\rho_A = \frac{1}{N} + \frac{w}{4N^2} \left\{ (3N^2 - 11N + 8)a + (N^2 + 3N - 8)b - (N^2 - 3N + 8)c - (3N^2 - 5N - 8)d \right\}
\]
\[
\rho_A \rho_B = 1 + \frac{w}{2} \left\{ (3N - 8)a + Nb - Nc - (3N - 8)d \right\}.
\]

(iii) Finally for IM updating we have
\[
\frac{\mu_i}{\lambda_i} = \begin{cases} 
\left[ (\gamma + \delta)/2\beta \right] \cdot \left[ (2\beta + \gamma + 3\delta)/(2\beta + 2\gamma + 2\delta) \right] & (i = 1) \\
\left[ (\gamma + \delta)/(\alpha + \beta) \right] \cdot \left[ (\alpha + \beta + \gamma + 3\delta)/(2\alpha + 2\beta + \gamma + \delta) \right] & (i = 2) \\
\left[ (\gamma + \delta)/(\alpha + \beta) \right] \cdot \left[ (\alpha + \beta + 2\gamma + 2\delta)/(3\alpha + \beta + \gamma + \delta) \right] & (3 \leq i \leq N - 3) \\
\left[ 2\gamma/(\alpha + \beta) \right] \cdot \left[ (2\alpha + 2\beta + 2\gamma)/(3\alpha + \beta + 2\gamma) \right] & (i = N - 2) \\
\left[ 2\gamma/(\alpha + \beta) \right] \cdot \left[ (2\alpha + 2\beta + 2\gamma)/(3\alpha + \beta + 2\gamma) \right] & (i = N - 1)
\end{cases}
\]

For strong selection, \(\rho_A > 0\) holds if and only if \([ (\gamma + \delta)/(\alpha + \beta) \cdot [(\alpha + \beta + \gamma + 3\delta)/(3\alpha + \beta + \gamma + \delta)] < 1\), which is equivalent to \((3a + b)(a + b) > (c + d)(c + 3d)\). For weak selection, we have
\[
\rho_A = \frac{1}{N} + \frac{w}{6N^2} \left\{ (5N^2 - 17N + 12)a + (3N^2 + 5N - 12)b \\
- (3N^2 - 5N + 12)c - (5N^2 - 7N - 12)d \right\}
\]
and
\[
\frac{\rho_A}{\rho_B} = 1 + \frac{w}{3} \left\{ (5N - 12)a + 3Nb - 3Nc - (5N - 12)d \right\}.
\]

If the interaction graph is complete, the total payoffs of \(A\) and \(B\) players are
\[
e_A(i) = (i - 1) \cdot a + (N - i) \cdot b
\]
\[
e_B(i) = i \cdot c + (N - i - 1) \cdot d.
\]
Substituting \(e_A(i)\) for \(a, b\) and \(e_B(i)\) for \(c, d\) leads to the results shown in the main text.

References


Ohtsuki, H. & Nowak, M. A. 2006 Fixation probabilities in evolutionary games on graphs. (in preparation)


Traulsen, A. & Claussen, J. C. 2004 Similarity-based cooperation and spatial segregation.


Figure captions

**Fig 1:** Evolutionary games on graphs (or social networks) are characterized by an interaction graph and a replacement graph. The interaction graph specifies who is interacting with whom in terms of the evolutionary game that is under consideration. The replacement graph specifies the local neighborhood for updating the strategies (that is who competes with whom in the selection process). The simplest geometry is that both the replacement graph and the interaction graph are given by the same cycle, where each individual has exactly two neighbors and there is no boundary. This one-dimensional geometry allows an exact calculation which can be compared with approximative methods for other geometries.

**Fig 2:** For ‘birth-death’ (BD) updating, only the payoff of the two individuals right at the boundary between two clusters matter. The cluster of $A$ players expands if $a+b > c+d$. For ‘death-birth’ (DB) and ‘imitation’ (IM) updating, the payoff of the four individuals closest to the boundary between two clusters matters. The cluster of $A$ players increases
if $2a + (a + b) > (c + d) + 2d$ for DB updating and if $2a + 3(a + b) > 3(c + d) + 2d$ for IM updating. All these results hold for weak selection. This figure gives the intuitive explanation why cooperation can be favored over defection for DB and IM, but not for BD updating.

**Short title for page headings:** games on cycles
Figure 1

Replacement graph

Interaction graph
a. Birth-death (BD) updating

b. Death-birth (DB) or imitation (IM) updating

\[ 2a \quad a + b \quad c + d \quad 2d \]